## Sur la topologie des courbes algébriques planes

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## Outline

(1) General algebraic problem
(2) Details of the algorithm
(3) Implementation and experiments

Luis Peñaranda (LORIA)

## Topology and some geometry of real algebraic plane curves

Input curve: $f(x, y)=0$ with $f \in \mathbb{Q}[x, y]$



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Curve: square free polynomial $f \in \mathbb{Q}[x, y]$.
A point $p=(\alpha, \beta) \in \mathbb{C}^{2}$ is $(x$ - $)$ critical if $f(p)=f_{y}(p)=0$

- singular if $f_{x}(\mathbf{p})=0$
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## Previous work

## Mainly 2 approaches

## Subdivision

- only guarantee the drawing up to some precision
- need to go up to the theoretical separation bound to be certified
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## Cylindrical Algebraic Decomposition-based with sub-resultant and lifting

Several variants:

- use Sturm-Habitch sequences or just principal S-H coefficients
- use generic position assumption
- use several projections
- shear and shear back


## CAD-based method

(1) Projection

Compute $x$-coordinates critical points: $\alpha_{i}$
(3) Lifting

Compute intersection points between the curve and the fiber $x=\alpha_{i}$
Compute with polynomial with algebraic coefficients
© Adjacencies
Count the number of branches : connected to the left and right

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## General idea

- Replace sub-resultant sequences and computations with algebraic coefficient polynomials by
Gröbner basis and Rational Univariate Representations
- Identify local topology at critical points using multiplicities and refinement
- Compute adjacencies with a vertical rectangular decomposition using multiplicities



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## Our algorithm

## Based on

- incremental work upon [WS05] and [CFPR08]
- Gröbner basis and Rational Univariate Representation of critical points.
[WS05] R. Seidel and N. Wolpert. On the Exact Computation of the Topology of Real Algebraic Curves. SoCG05. [CFPR08] F. Cazals, J.-C. Faugère, M. Pouget, and F. Rouillier. Ridges and Umbilics of Polynomial Parametric Surfaces, in Geometric Modeling and Algebraic Geometry, Springer.


## Specifications

- compute the exact topology (output a straight line graph)
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## Algorithm outline

(1) Compute isolating boxes for critical points, easily refinable with the RUR.
(2) Topology at extreme points:

(3) Topology at singular points:

(9) Topology in non critical cells of the induced vertical rectangular decomposition of the plane.

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Rational Univariate Representation

Univariate real solving

## Algebraic tools

- Univariate root isolation for polynomials with rational coefficients: Descartes algorithm
- Solve zero-dimensional systems with Rational Univariate Representation (RUR) preserves
(1) real roots
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## Topology at extreme points

(1) Isolate the extreme system

$$
I_{e}=\mathbb{I}\left(f, f_{y}, f_{x} \neq 0\right)=\mathbb{I}\left(f, f_{y}, T f_{x}-1\right) \cap \mathbb{Q}[x, y]
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## Topology at singularities

(1) Isolate singular points in boxes
(2) Compute multiplicities $k$ in fibers
(3) Refine the box to avoid the curve $f_{y^{k}}=\frac{\partial^{k} f}{\partial y^{k}}$
(9) Refine the box to avoid top/bottom crossings


## Rectangle decomposition of the plane

- the topology is known inside critical boxes
- compute a vertical decomposition of the plane wrt these boxes
- compute intersections of the curve with the decomposition



## Greedy connection algorithm using multiplicities

Overlapping of extreme point boxes: need parity of multiplicity in fiber


## Complexity analysis

## Theorem

The algorithm runs in $\widetilde{\mathcal{O}}_{\mathrm{B}}\left(\mathrm{Rd}^{4}\left(\mathrm{~d} \mathrm{\tau s}+\mathrm{s}^{2}\right)\right)$, where

- R: number of real critical points,
- d : degree of the polynomial f ,
- $\tau$ : maximum coefficient bitsize of f ,
- s: maximum bitsize of
- the separation bound of $\mathrm{I}_{\mathrm{c}}$, and
- the distance between a singular point and its isolating curve (worst case $s=d^{3} \tau$ ).


## Implementation

Isotop: 7638 lines of Maple code using packages:

- FGb for Gröbner basis (Faugère)
- RS for RUR and isolation (Rouillier)
- faster on non generic and high degree curves
- robust
- exact
http://www.loria.fr/equipes/vegas/isotop/


## Experiments

We ran large-scale tests, comparing to Brown's Cad2d, MPII's Alcix, González Vega's Top and Wolpert's Insulate.

| example index | $\mathrm{r}=\frac{\mathrm{time}_{\text {cad2d }}}{\text { time }{ }_{\text {Isoror }}}$ | example index | $r=\frac{t i m e e_{\text {Alcix }}}{\text { time }{ }_{\text {Isorop }}}$ |
| :---: | :---: | :---: | :---: |
| $1 \ldots 114$ | $1 / 43 \leq r<1 / 3$ | $1 \ldots 17$ | $1 / 5<\mathrm{r}<1 / 3$ |
| 115... 382 | $1 / 3 \leq r \leq 1$ | $18 \ldots 33$ | $1 / 3 \leq r \leq 1$ |
| $383 \ldots 479$ | $1<\mathrm{r} \leq 3$ | $34 \ldots 290$ | $1<\mathrm{r} \leq 3$ |
| $480 \ldots 513$ | $3 \leq r \leq 160$ | 291... 465 | $3 \leq r \leq 64$ |
| $514 \ldots 585$ | Cad2d timeout | $466 \ldots 585$ | Alcix timeout |

## Conclusion

Summary of our contribution:

- deals with non generic case,
- gives results in the original coordinate system (identifies vertical asymptotes),
- avoids Sturm Habitch (principal coefficient),
- uses RUR,
- no restriction on singularities,
- enhancement for extreme points and the connection algorithm,
- output sensitive complexity analysis.

